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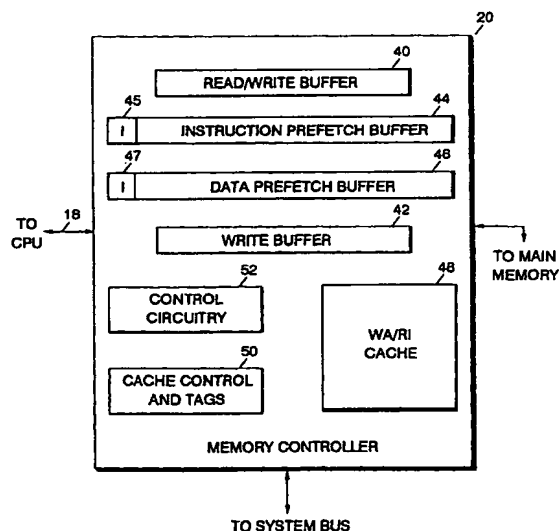
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(54) **Memory controller**

(57) An improved memory controller (20) within a data processing system (10) having a look-aside cache architecture is disclosed. The data processing system (10) includes a processor (12) having an upper level cache (14) associated therewith, a memory controller (20) having an associated controller memory (48), a processor bus (18) coupled between the processor and the memory controller (20), and a main memory (22). The data processing system (10) further includes a lower level cache (16) coupled to the processor bus (18) in parallel with the processor (12) and memory controller (20). According to a first aspect of the present invention, the memory controller (20) includes logic, which in response to receipt of a write request that will not be serviced by the lower level cache (16) and for which the associated data is not a replaced modified cache line, stores the associated data within the controller memory (48) associated with the memory controller (20), thereby optimizing data storage within the data processing system (10). According to a second aspect of the present invention, the memory controller (20) includes logic, which in response to receipt of a request for information residing only in main memory (22), fetches the requested information from main memory (22) and stores additional information adjacent to said requested data in main memory (22) within a prefetch buffer (44, 46), thereby minimizing access time to the prefetched information.

FIG. 2



EP 0 800 137 A1

## Description

### Technical Field

The present invention relates in general to a data processing system and in particular to an apparatus for managing data storage within a data processing system. Still more particularly, the present invention relates to a memory controller within a data processing system having a look-aside cache architecture which caches stack operations and prefetches selected information for possible subsequent access by the processor.

### Description of the Related Art

To decrease latency, data processing systems increasingly utilize some configuration of cache memory. As is well-known to those skilled in the art, a cache is a small amount of fast, expensive, zero wait state memory utilized to store a copy of frequently accessed instructions and data residing in main memory. The latest generation of personal computers, which utilize 80486, Intel Pentium, IBM PowerPC, or similar processors typically include an on-chip level one (L1) processor cache. (Pentium is a trademark of Intel Corp. and IBM and Power PC are trademarks of IBM Corp.). In addition, these personal computers frequently include a level two (L2) cache to further enhance system performance. Cache systems having both L1 and L2 caches are typically configured in one of two ways. In the first cache system configuration, the L2 cache is interfaced in a serial fashion between the processor and the system or memory bus. In this configuration, commonly referred to as a look-through or in-line configuration, the processor cannot communicate directly with the memory or system bus, but communicates through the interface provided by the L2 cache controller.

Although an in-line L2 cache configuration generally provides optimal performance, many personal computer systems are designed to support optional L2 caches in a look-aside configuration in order to lower the price of an entry-level computer system while providing the option to install an L2 cache to improve performance. In a look-aside configuration, the L2 cache is coupled to the processor bus in parallel with both the processor and the memory controller and may therefore conveniently be mounted on a pluggable module connected with the processor bus.

In computer systems which utilize a look-aside L2 cache configuration, the L2 cache and memory controller each begin a processor memory read cycle simultaneously in response to the processor initiating a memory read. In response to an L2 cache read hit, the L2 cache signals the memory controller to abort the indicated memory read and returns the requested data to the processor in zero wait states. However, in the event of an L2 cache read miss, the memory controller fetches the requested data from main memory and returns the

data to the processor as if the L2 cache were not present. Since the L2 cache and the memory controller both begin to service a processor data read request simultaneously, a computer system having a look-aside cache architecture incurs no added penalty for an L2 cache miss during a data read.

Implementing an L2 cache within a computer system utilizing a look-aside configuration typically has a concomitant performance penalty, however. For example, in the case of a cache miss of a look-aside L2 cache during a data write, a performance penalty is incurred since the L2 cache cannot obtain control of the processor bus in order to fetch the requisite cache line while the processor is writing the data to the memory controller. Consequently, look-aside L2 caches typically do not implement cache line allocation on write misses. In addition, contention for the processor bus also reduces system performance during I/O operations because the processor cannot access the L2 cache during an I/O operation. A further limitation of a look-aside L2 cache configuration is that it does not efficiently support cache line sizes larger than the L1 cache line size. In contrast, in-line L2 cache lines are frequently designed to be twice the length of L1 cache lines in order to reduce cache miss ratios by prefetching instructions and data based upon the statistical probability of data locality.

As should thus be apparent, it would be desirable to provide an improved method and system for implementing an optional look-aside L2 cache within a data processing system. In particular, it would be desirable to provide an improved cache system within a data processing system having a look-aside L2 cache configuration which support allocation on L2 write misses and which enable the prefetching of data and instructions.

### DISCLOSURE OF THE INVENTION

An improved memory controller within a data processing system having a look-aside cache architecture is disclosed. The data processing system includes a processor having an upper level cache associated therewith, a memory controller having an associated controller memory, a processor bus coupled between the processor and the memory controller, and a main memory. The data processing system further includes a lower level cache coupled to the processor bus in parallel with the processor and memory controller. According to a first aspect of the present invention, the memory controller includes logic, which in response to receipt of a write request that will not be serviced by the lower level cache and for which the associated data is not a replaced modified cache line, stores the associated data within the controller memory associated with the memory controller, thereby optimizing data storage within the data processing system. According to a second aspect of the present invention, the memory controller includes logic, which in response to receipt of a request for infor-

mation residing only in main memory, fetches the requested information from main memory and stores additional information adjacent to said requested data in main memory within a prefetch buffer, thereby minimizing access time to the prefetched information.

## BRIEF DESCRIPTION OF THE DRAWINGS

The invention will now be described by way of example only, with reference to the accompanying drawings, in which:

**Figure 1** illustrates a high-level block diagram of a data processing system in accordance with the method and system of the present invention;

**Figure 2** depicts a more detailed block diagram of a memory controller in accordance with the method and system of the present invention;

**Figure 3** illustrates a high-level logic flowchart of a preferred embodiment of the method of the present invention;

**Figures 4A and 4B** are a high-level logic flowchart of a preferred embodiment of the method utilized by a memory controller which employs the present invention to service an instruction fetch request;

**Figures 5A and 5B** are a high-level logic flowchart of a preferred embodiment of the method utilized by a memory controller which employs the present invention to service a data write request; and

**Figures 6A and 6B** are a high-level logic flowchart of a preferred embodiment of the method utilized by a memory controller which employs the present invention to service a data read request.

## DETAILED DESCRIPTION OF THE INVENTION

With reference now to the figures and in particular with reference to **Figure 1**, there is illustrated a block diagram of a preferred embodiment of a data processing system in accordance with the method and system of the present invention. As will be appreciated by those skilled in the art, many of the details of data processing system **10** that are not relevant to the present invention have been omitted for the purpose of clarity. As illustrated, data processing system **10** includes a central processing unit (CPU) **12** which executes software instructions. While any appropriate microprocessor can be utilized for CPU **12**, CPU **12** is preferably one of the PowerPC line of microprocessors available from IBM Microelectronics. Alternatively, CPU **12** can be implemented as an Intel Pentium or an 80486 microprocessor. To improve data and instruction access times, CPU **12** is equipped with an on-board level one (L1) cache

**14**. Although in the following description the cache line size of L1 cache **14** is described as being x bytes in length, in a preferred embodiment of the present invention in which the word length of CPU **12** is 8 bytes, the cache line length of L1 cache **14** is 32 bytes. CPU **12** is coupled to processor bus **18**, which preferably has a bandwidth of 8 bytes, to facilitate communication of data and instructions between CPU **12**, L2 cache **16** and memory controller **20**.

As depicted, L2 cache **16** is coupled to processor bus **18** in parallel with CPU **12** and memory controller **20** in a look-aside cache configuration. Accordingly, read and write requests transmitted by CPU **12** via processor bus **18** are received concurrently by memory controller **20** and L2 cache **16**. In response to an L2 cache hit, L2 cache **16** signals memory controller **20** to abort the indicated operation and returns the requested data to CPU **12** in zero wait states. L2 cache **16** preferably has a cache line length of X bytes to avoid the complications inherent in supporting multiple caches having diverse cache line sizes on a shared bus. As illustrated, L2 cache **16** includes an L2 cache controller **17**, which controls the operation of L2 cache **16**. Thus, L2 cache controller **17** maintains L2 cache coherency by enforcing a selected coherency protocol, determines whether data associated with memory addresses within main memory **22** are cacheable, or capable of residing within L2 cache **16**, and performs many other conventional cache management functions.

Data processing system **10** further includes memory controller **20**. Memory controller **20** contains logic circuitry which fetches data and instructions from main memory **22** in response to receipt of read and write requests from CPU **12** which cannot be serviced by L2 cache **16**. Thus, memory controller **20** provides a memory interface between CPU **12** and main memory **22**. In addition, memory controller **20** includes logic circuitry which provides a system bus interface between system bus **24** and CPU **12** and main memory **22**. In a preferred embodiment of the present invention, the system bus interface within memory controller **20** supports memory mapped I/O by transmitting data received from CPU **12** to system bus **24** if the specified address maps to an address assigned to an I/O device.

As is further illustrated within Figure 1, data processing system **10** includes read only memory (ROM) **26**, I/O adapter **28**, secondary storage **30**, and display adapter **32**, which are each coupled to system bus **24**. ROM **26** and secondary storage **30** provide storage for operating system and application programs and data. I/O adapter **28** supports the attachment of input devices, such as a mouse and keyboard, to data processing system **10** to enable a user to input data and instructions. Display adapter **32** enables the attachment of a video display device to output data to a user.

Referring now to **Figure 2**, there is depicted a more detailed pictorial representation of the logical structure of memory controller **20** in accordance with the method

and system of the present invention. As illustrated, memory controller 20 contains a conventional read/write buffer 40 and write buffer 42. Read/write buffer 40 is utilized to buffer data transmitted to and received from CPU 12 via processor bus 18. Write buffer 42 is utilized to buffer data to be written to main memory 22. Each of read/write buffer 40 and write buffer 42 preferably has the same length as a cache line of L1 cache 14 in order to support efficient data transfers, for example, burst transfers between memory controller 20 and CPU 12.

In accordance with the present invention, memory controller 20 further includes an instruction prefetch buffer (IPB) 44 and a data prefetch buffer (DPB) 46. IPB 44 and DPB 46 are utilized by memory controller 20 to prefetch data and instructions for CPU 12. As described above, based upon the principle of locality of reference, it has been shown that cache miss ratios are greatly reduced by implementing a 2:1 L2 to L1 cache line size ratio in order to prefetch an additional L1 cache line of data and instructions during each fetch from memory. Because diverse L2 and L1 cache line sizes are not easily supported when a look-aside cache configuration is utilized, memory controller 20 fetches two cache lines of data or instructions from main memory 22 during particular fetch operations and stores the data or instructions contained within the cache line not immediately requested by CPU 12 within the appropriate one of IPB 44 and DPB 46. Thus, as will be described in greater detail below, memory controller 20 supports the prefetching of data and instructions in conjunction with a look-aside configuration of L2 cache 16.

According to another aspect of the present invention, memory controller 20 also includes write allocate/read invalidate (WA/RI) cache 48 and its associated cache control and tags 50. Within conventional data processing systems which implement a look-aside L2 cache, the memory controller simply writes data received from the processor to the main memory in response to an L2 cache write miss. Thus, a conventional look-aside cache typically does not allocate a cache line in response to a write miss. This storage management policy is beneficial if the data to be written is a replaced L1 or L2 cache line since the probability that the replaced cache line will soon be accessed again is small. However, if the data write is a stack operation, failure to allocate a cache line in response to a write miss degrades system performance.

As is well known to those skilled in the art, a stack is a logical first-in/last-out (FILO) queue which is utilized to save parameters during procedure calls and other software operations which save parameters. Stack operations tend to write parameters to a data location first (a "push") and thereafter read the data location (a "pop"). Since the stack data will typically be read only once, stack data is considered invalid following a pop. According to the present invention, in order to efficiently support push stack operations, WA/RI cache 48 within memory controller 20 allocates a cache line on write

misses of L2 cache 16 that are single word (non-burst) writes. WA/RI cache 48 does not allocate a cache line on multiple-word writes (burst writes) since burst writes typically represent replaced cache lines that no longer need to be cached. In addition, WA/RI cache 48 invalidates data following a read hit (a pop).

Finally, memory controller 20 includes control circuitry 52, which manages the operation of memory controller 20 in accordance with the logical process illustrated within Figures 3-6. Upon review of Figures 3-6, those skilled in the art will appreciate that many operations depicted in a serial fashion therein may in practice be performed in parallel. With reference first to Figure 3, there is illustrated a high-level logic flowchart of the operation of memory controller 20 in accordance with the method and system of the present invention. As illustrated, the process begins at block 60 and thereafter proceeds to block 62, which illustrates a determination of whether or not an operation request received from CPU 12 via processor bus 18 is an instruction fetch request. In response to a determination that the operation request is not an instruction fetch request, the process passes to block 64. However, in response to a determination that the operation request is an instruction fetch request, the process proceeds through off-page connector A to on-page connector A of Figure 4.

Referring now to Figure 4, there is depicted a high-level logic block diagram of a preferred embodiment of the process utilized by memory controller 20 to prefetch instructions in accordance with the method and system of the present invention. As illustrated, the process proceeds from on-page connector A to block 70, which depicts a determination of whether or not the instruction fetch request resulted in an L2 cache hit. If L2 cache 16 stores the requested instructions, L2 cache 16 signals memory controller 20 to abort its operation. Therefore, if the instructions associated with a specified memory address are stored within L2 cache 16, the process proceeds from block 70 to block 118 and terminates. However, if a determination is made at block 70 that the instruction fetch request resulted in a L2 cache miss, L2 cache 16 cannot service the instruction fetch request and the process passes to block 72.

Block 72 depicts a determination of whether or not the instructions specified by the instruction fetch request are stored within WA/RI cache 48. If not, the process passes from block 72 to block 80. However, if a determination is made at block 72 that WA/RI cache 48 stores the requested instructions, the process proceeds from block 72 to blocks 74-78, which illustrate memory controller 20 transmitting the requested instructions to CPU 12 via processor bus 18, writing back the WA/RI cache line containing the requested instructions to main memory 22, and invalidating the WA/RI cache line containing the requested instructions. Thereafter, the process passes to block 118 and terminates. Returning to block 72, if a determination is made that the requested instructions are not stored within WA/RI cache 48, the process

passes to block 80, which illustrates a determination of whether or not the requested instructions are stored within DPB 46. Although the operation request issued by CPU 12 is an instruction fetch request, memory controller 20 determines whether DPB 46 stores the requested instructions since computer architectures typically permit information to be accessed as instructions or data in order to support self-modifying code. In response to a determination that the requested instructions are stored within DPB 46, the process passes from block 80 to block 82, which illustrates memory controller 20 transmitting the requested instructions to CPU 12. Next, the process proceeds to blocks 84-86, which illustrate invalidating DPB 46 by setting invalid bit 47 if a full L1 cache line was transmitted to CPU 12. The process then passes to block 118 and terminates.

Returning to block 80, if a determination is made that the requested instructions are not stored within DPB 46, the process passes to block 88, which depicts a determination of whether or not the requested instructions are stored within IPB 44. In response to a determination that the requested instructions are stored within IPB 44, the process proceeds to block 90, which illustrates memory controller 20 transmitting the requested instructions to CPU 12. Next, the process passes to block 92, which depicts determining whether or not a full L1 cache line of instructions was transmitted to CPU 12. If not, the process simply passes to block 118 and terminates. However, if a full L1 cache line was transmitted, the process proceeds to block 94, which illustrates memory controller 20 invalidating the contents of IPB 44 by setting invalid bit 45. The process then proceeds to block 96, which depicts a determination of whether or not the x bytes (x is the cache line length of L1 cache 14) that follow the requested instructions within main memory 22 are cacheable. If not, the process passes to block 118 and terminates. However, if a determination is made that the next x bytes within main memory 22 are cacheable, the process proceeds to block 98, which illustrates memory controller 20 fetching the x bytes following the requested instructions, storing them within IPB 44, and clearing invalid bit 45. Thereafter, the process passes to block 118 and terminates.

Returning to block 88, if a determination is made that the requested instructions do not reside within IPB 44, the process passes to block 100, which depicts a determination of whether or not the requested instructions represent a full L1 cache line and whether or not both the addresses containing the requested instructions and the following x bytes of information are both cacheable. If so, the process proceeds to block 102, which depicts fetching two L1 cache line lengths of bytes of information from main memory 22. Then, as illustrated at block 104, memory controller 20 transmits the first x bytes of instructions to CPU 12 and stores the second x bytes within IPB 44. Thus, memory controller 20 effectively prefetches a second L1 cache line of instructions because of the likelihood of a subsequent request for

instructions within the second x bytes of information. The process then passes to block 118 and terminates.

Returning to block 100, if a determination is made that either a full L1 cache line was not requested by CPU 12 or that 2X bytes are not cacheable, the process passes to block 106, which illustrates a determination of whether or not the x bytes within main memory 22 which contain the address of the requested instruction(s) are cacheable. If not, the process passes to block 108, which depicts memory controller 20 fetching the requested instruction(s) from main memory 22 and sending the requested instructions to CPU 12. The process then passes to block 118 and terminates. However, if a determination is made at block 106 that x bytes of information containing the requested instructions are cacheable, the process passes to block 110 and 112, which illustrate memory controller 20 fetching the X bytes containing the requested instructions from main memory 22 and transmitting the requested instructions to CPU 12. Next, a determination is made at block 114 whether or not x bytes, which comprise a full L1 cache line, were sent to CPU 12. If so, the process passes to block 118 and terminates. However, if less than a full cache line of instructions was sent to CPU 12, the process passes to block 116, which depicts storing the X fetched bytes of information within IPB 44 and marking them valid by clearing invalid bit 45. Thereafter, the process passes to block 118 and terminates.

Referring again to Figure 3, if a determination is made at block 62 that the CPU operation request received by memory controller 20 is not an instruction fetch request, the process passes to block 64, which depicts a determination of whether or not the CPU operation request is a data write request. If so, the process proceeds from block 64 through off-page connector B to Figure 5, which illustrates a preferred embodiment of the process utilized by memory controller 20 to service data write requests.

With reference now to Figure 5, the process utilized by memory controller 20 to service data write requests begins at on-page connector B and thereafter proceeds to block 130, which illustrates a determination of whether or not the data write request will be serviced by L2 cache 16. As described above with reference to the instruction fetch request, L2 cache 16 signals that a copy of the data stored at the specified address resides within L2 cache 16 by transmitting an abort signal to memory controller 20. In response to receipt of the abort signal indicating that the data write request will be serviced by L2 cache 16, the process proceeds from block 130 to block 172, where the process terminates. However, in response to a determination that the data write request will not be serviced by L2 cache 16, the process proceeds from block 130 to block 132, which illustrates a determination of whether or not the data associated with the data write request is a cache line cast out of (replaced from) L1 cache 14 or L2 cache 16 or is locked or is otherwise noncacheable. If so, the process proceeds

from block 132 to block 134, which depicts memory controller 20 writing the data associated with the data write request to the specified address within main memory 22. Next, as depicted at block 136, memory controller 20 snoops WA/RI cache 48, IPB 44, and DPB 46 and invalidates any data within memory controller 20 corresponding to the specified address. The process then passes to block 172 and terminates. Returning to block 132, if a determination is made that the data associated with the data write request is not a cache line cast out of L1 cache 14 or L2 cache 16 or locked or noncacheable, the process passes to block 138, which depicts determining whether or not DPB 46 stores data corresponding to the specified addresses. If so, the process proceeds from block 138 to block 140, which illustrates merging the data associated with the data write request with the data stored within DPB 46. Next, as illustrated at block 142, the information within DPB 46 is written into WA/RI cache 48. Thereafter, the contents of DPB 46 are invalidated by setting invalid bit 47 and the process passes to block 172, where the process terminates.

Returning to block 138, if a determination is made that DPB 46 does not contain data associated with the specified address, the process proceeds from block 138 to block 146, which illustrates a determination of whether or not IPB 44 stores information associated with the specified addresses. If so, the process proceeds to blocks 148-152, which like block 140-144, depict memory controller 20 merging the data associated with the data write request with the contents of IPB 44, storing the content of IPB 44 within WA/RI cache 48, and thereafter invalidating IPB 44 by setting invalid bit 45. The process then passes to block 172 and terminates. Returning to block 146, if a determination is made that IPB 44 does not store information associated with a specified address, the process proceeds from block 146 to block 154, which illustrates a determination of whether or not information associated with the specified address is stored within WA/RI cache 48. The determination illustrated at block 154 is preferably made by comparing selected bits within the specified address with address tags stored within cache control and tags 50. If the selected bits within the specified address match one of the address tags stored within cache control and tags 50, indicating that WA/RI cache 48 stores information associated with the specified address, the process passes from block 154 to block 156, which illustrates memory controller 20 updating a WA/RI cache line with the data associated with the data write request. The process then passes to block 172 and terminates.

Returning to block 154, if a determination is made that the data write request results in a cache miss of WA/RI cache 48, the process proceeds from block 154 to block 158, which illustrates allocating a cache line within WA/RI cache 48 for the data associated with the data write request. Next, as depicted at block 160, memory controller 20 fetches X bytes of data containing the specified address from main memory 22 and stores the

5 fetched data within read/write buffer 40. In addition, memory controller merges the data associated with the data write request with the contents of read/write buffer 40. The process then proceeds to block 162, which illustrates a determination of whether or not the replaced WA/RI cache line has been modified. For example, the determination depicted at block 162 may be made by examining the coherency protocol bit associated with the cache line. If the cache line is marked as dirty, the process proceeds to block 164, which illustrates writing the replaced WA/RI cache line to main memory 22. The process then proceeds from either block 164 or block 162 to block 168, which depicts storing the contents of read/write buffer 40 into the allocated WA/RI cache line. The cache line is then marked as modified (valid) as illustrated at block 170. Thereafter, the process passes to block 172 and terminates.

Referring again to Figure 3, if a determination is made at block 64 that the CPU operation request received at memory controller 20 is not a data write request, the process passes to block 66, which depicts a determination of whether or not the CPU operation request is a data read request. If not, the process passes to block 68 and terminates. However, if a determination is made at block 66 that the CPU operation request is a data read request, the process proceeds through off-page connector C to on-page connector C of Figure 6. Referring now to Figure 6, there is illustrated a high-level flowchart of a preferred embodiment of the method utilized by the present invention to service a data read request. As illustrated, the process passes from on-page connector C to block 180, which illustrates a determination of whether or not the CPU operation request will be serviced by L2 cache 16. If so, the process simply passes to block 232 and terminates. However, if a determination is made at block 180 that L2 cache 16 will not service the CPU operation request, the process passes to block 182, which illustrates a determination of whether or not data associated with the address specified within the data read request is stored within WA/RI cache 48. If so, the process proceeds from block 182 to blocks 184-190, which depict the read invalidate operation of WA/RI cache 48. First, as illustrated at block 184, the data associated with the address specified within the data read request is transmitted to CPU 12 via processor bus 18. Next, the process passes to block 186, which illustrates a determination of whether or not X bytes of data, a full L1 cache line, were transmitted to CPU 12. If not, the process passes to block 232 and terminates. However, if a full L1 cache line was transmitted to CPU 12, the process proceeds to block 188-190, which depict memory controller 20 writing back the WA/RI cache line containing the requested data to main memory 22 and marking the cache line invalid. The process then passes to block 232 and terminates.

Returning to block 182, if a determination is made that the requested data is not stored within WA/RI cache

48, the process proceeds to block 192, which depicts a determination of whether or not the requested data is stored within IPB 44. In response to a determination that the requested information is stored within IPB 44, the process proceeds to block 194, which illustrates returning the requested data to CPU 12. Then, as depicted at block 196, a determination is made whether or not the requested data comprised a full L1 cache line. If not, the process passes to block 232 and terminates. However, if a determination is made that the requested data comprised a full L1 cache line, the process proceeds to block 198, which illustrates invalidating the contents of IPB 44 by setting invalid bit 45. The process then passes to block 232 and terminates. Returning to block 192, if a determination is made that IPB 44 does not contain the requested data, the process passes to block 200, which depicts a determination of whether or not the requested data resides within DPB 46. If so, the process proceeds to block 200 to block 202, which illustrates returning the requested data to CPU 12. Next, as depicted at block 204, a determination is made whether or not the requested data comprised a full L1 cache line. If not, the process passes to block 232 and terminates. However, if the requested data comprised a full L1 cache line, the process proceeds to block 206, which illustrates memory controller 20 invalidating the contents of DPB 46 by setting invalid bit 47. The process proceeds from block 206 to block 208, which depicts a determination of whether or not the X bytes within main memory 22 which follow the X bytes of requested data are cacheable. If not, the process passes to block 232 and terminates. However, in response to a determination that the next X bytes of information within main memory 22 are cacheable, the process passes from block 208 to block 210, which depicts memory controller 20 fetching the subsequent X bytes of information from main memory 22 and storing them within DPB 46. In addition, memory controller 20 marks DPB as valid by clearing invalid bit 47. Block 210 again illustrates memory controller 20 prefetching data based upon the principle of locality of reference in order to potentially avert future main memory accesses which result from L2 cache misses.

Returning to block 200, if a determination is made that the requested data does not reside within DPB 46, the requested data must be fetched from main memory 22 and the process passes to block 212. Block 212 depicts a determination of whether or not the data read request requests X bytes of information and whether or not the 2X bytes of information within main memory 22 containing the specified address are cacheable. If so, the process proceeds from block 212 to block 214, which illustrates fetching the 2X bytes of information containing the specified address from main memory 22. Then, as depicted at blocks 216-218, memory controller 20 transmits the first X bytes of information to CPU 12 and stores the second X bytes of information within DPB 46, marking them valid by clearing invalid bit 47. Thereafter, the process terminates at block 232. Returning to

block 212, if a determination is made that the requested data does not comprise a full L1 cache line or that two cache lines of data are not cacheable, the process passes to block 220, which depicts determining whether or not the X bytes of information following the specified address within main memory are cacheable. If so, the process proceeds from block 220 to block 222, which depicts fetching the X bytes of data following the specified address from main memory 22 and sending the requested data to CPU 12. Next, as illustrated at block 224, a determination is made whether or not the requested data comprises a full L1 cache line. If so, the process passes to block 232 and terminates. However, if the requested data does not comprise a full L1 cache line, the process proceeds to block 226, which illustrates memory controller 20 storing the X bytes of data fetched from main memory 22 within DPB 46 and clearing invalid bit 47. The process then passes to block 232 and terminates.

Returning to block 220, if a determination is made that the X bytes of data within main memory 22 containing the specified address are not cacheable, the process proceeds to blocks 228-230, which depicts memory controller 20 fetching only the requested data from main memory 22 and transmitting the requested data to CPU 12. Thereafter, the process terminates at block 232.

As should thus be apparent, the present invention provides an improved method and system for managing the storage of data within a data processing system having a look-aside cache configuration. In particular, the present invention optimizes data access times by providing a write-allocate/read-invalidate (WA/RI) cache within the memory controller in order to efficiently handle stack operations. Furthermore, according to the present invention, the memory controller includes prefetch buffering in order to minimize the latency incurred by L2 look-aside cache misses.

## Claims

1. A memory controller (20) for managing storage of data within a data processing system (10) having a look aside cache configuration, said data processing system including a processor (12) having a upper level cache (14) associated therewith, a controller memory (43) for coupling to said memory controller, a processor bus (18) for coupling between said processor and said memory controller, a lower level cache (16) coupled to said processor bus in parallel with said processor, and a main memory (22), wherein said upper level cache and said lower level cache each include one or more cache lines, said memory controller comprising:

means, responsive to receipt at said memory controller of a write request and associated data for a specified address within said main memory, for determining if said write request

- will be serviced by said lower level cache and if said associated data is a modified cache line replaced from either said upper level cache or said lower level cache;  
 means, responsive to a determination that said write request will not be serviced by said lower level cache and that said associated data is a modified cache line replaced from either said upper level cache or said lower level cache, for storing said associated data at said specified address within said main memory; and  
 means, responsive to a determination that said write request will not be serviced by said lower level cache and that said associated data is not a modified cache line replaced from either said upper level cache or said lower level cache, for storing said associated data within said controller memory associated with said memory controller, wherein data storage within said data processing system is optimized.
2. A memory controller (20) as claimed in Claim 1, and further comprising:  
 means, responsive to an access of said associated data by said processor (12), for invalidating said associated data within said controller memory (48).
3. A memory controller (20) as claimed in Claim 1, said controller memory (48) comprising an on-board cache memory within said memory controller.
4. A data processing system (10), comprising:  
 a processor (12);  
 a processor bus (18) coupled to said processor;  
 an upper level cache (14) coupled to said processor;  
 a lower level cache (16) coupled to said processor bus in parallel with said processor;  
 a main memory (22);  
 a memory controller (20) as claimed in any one of claim 1 to claim 3;  
 a system bus (24) coupled to said memory controller; and  
 one or more adapters (28, 32) coupled to said system bus for receiving inputs to said data processing system and presenting outputs of said data processing system to a user.
5. A memory controller (20) for use within a data processing system (10) including a processor (12) having a distributed cache memory (14, 16) associated therewith and a main memory, said memory controller comprising:  
 a prefetch buffer (44, 46);  
 means, responsive to a request by said processor for information, for determining if said requested information is stored within said distributed cache memory;  
 means, responsive to a determination that said requested information is not stored within said distributed cache memory, for determining whether or not said requested information is stored within said prefetch buffer within said memory controller;  
 means, responsive to a determination that said requested information is stored within said prefetch buffer, for transmitting said requested information to said processor; and  
 means, responsive to a determination that said requested information is not stored within said prefetch buffer, for fetching said requested information from said main memory for said processor and for storing additional information adjacent to said requested data in said main memory within said prefetch buffer, wherein access time of said processor to prefetched information is minimized.
6. A memory controller (20) as claimed in Claim 5, wherein said requested information comprises at least one instruction.
7. A memory controller as claimed in Claim 6, said memory controller further including a memory controller cache (48), wherein said means for determining whether or not said requested information is stored within distributed cache memory comprises means for determining whether or not said at least one instruction is stored within said memory controller cache.
8. A memory controller (20) as claimed in Claim 7, and further comprising:  
 responsive to a determination that said requested information is stored within said memory controller cache (48);  
 means for transmitting said at least one instruction to said processor;  
 means for storing a line of said memory controller cache which contains said at least one instruction within said main memory (22); and  
 means for invalidating said line of said memory controller cache.
9. A memory controller (20) as claimed in Claim 5, wherein said requested information comprises data.
10. A memory controller (20) as claimed in Claim 9, said memory controller further comprising a memory controller cache (48), wherein said means for determining whether or not said requested information



is stored within distributed cache memory comprises means for determining whether or not said at least one instruction is stored within said memory controller cache.

11. A memory controller (20) as claimed in Claim 6 or Claim 10, wherein said distributed cache memory includes at least an upper level cache (14) including one or more of cache lines having a cache line length of X bytes, said memory controller further comprising:

responsive to a determination that said requested data is stored within said memory controller cache:

means for transmitting said requested data to said processor (12);

means for determining whether or not said requested data comprises X bytes of data; and

means, responsive to a determination that said requested data comprises X bytes of data, for storing a line of said memory controller cache which contains said requested data within said main memory and for invalidating said line of said memory controller cache.

12. A memory controller (20) as claimed in Claim 6 or Claim 9, wherein said distributed cache memory includes at least an upper level cache (14) including one or more of cache lines having a cache line length of X bytes, said memory controller having an instruction prefetch buffer (44) and a data prefetch buffer (46), said memory controller further comprising:

means, responsive to a determination is that said requested information is stored within said data prefetch buffer, for invalidating information stored within said instruction prefetch buffer following said transmission of said requested information to said processor if said requested information comprises X bytes;

means, responsive to said invalidation of said information within said data prefetch buffer, for determining whether or not X bytes of information within said main memory adjacent to said X bytes of requested information are cacheable; and

means, responsive to a determination that X bytes of information within said main memory (22) adjacent to said X bytes of requested information are cacheable, for fetching from said main memory said X bytes of information adjacent to said X bytes of requested information and storing said X fetched bytes of information within said data prefetch buffer.

13. A memory controller (20) as claimed in Claim 9,

wherein said distributed cache memory (48) includes at least an upper level cache (14) including one or more of cache lines having a cache line length of X bytes, and wherein said means for fetching said requested information from said main memory (22) for said processor (12) and for storing additional information adjacent to said requested data in said main memory within said prefetch buffer comprises:

means for determining if said requested information comprises X bytes and if a following X bytes of information are cacheable;

responsive to a determination that said requested information comprises X bytes and a following X bytes of information are cacheable:

means for fetching said X bytes of requested information and said following X bytes of information from said main memory;

means for transmitting said X bytes of requested information to said processor;

means for storing said following X bytes of information within said prefetch buffer (44, 46);

means, responsive to a determination that said requested information does not comprise X bytes or that said following X bytes of information are not cacheable, for determining if X bytes of information within said main memory which contain said requested information are cacheable;

responsive to a determination that said X bytes of information within said main memory which contain said requested information are cacheable:

means for fetching from said main memory said X bytes of information which contain said requested information and for storing said X bytes of information fetched from said main memory within said prefetch buffer, if said requested information comprises less than X bytes of information; and

means for transmitting said requested information to said processor.

14. A data processing system (10), comprising:

a processor bus (18);

a processor (12) coupled to said processor bus;

a distributed cache memory (14, 16) coupled to said processor;

a main memory (22);

a memory controller (20) as claimed in Claim 5 coupled to said processor bus and to said main memory;

means, responsive to a determination that said requested information is not stored within said prefetch buffer (44, 46), for fetching said requested information from said main memory for

said processor and for storing additional information adjacent to said requested data in said main memory within said prefetch buffer; a system bus (24) coupled to said memory controller; and one or more adapters (28, 32) coupled to said system bus for receiving inputs to said data processing system and presenting outputs of said data processing system to a user.

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15. A data processing system (10) as claimed in Claim 14, wherein said distributed cache memory (14, 16) includes an upper level cache (14) coupled to said processor (12) and a lower level cache (16) coupled to said processor bus (18) in parallel with said processor and said memory controller (20).

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16. A data processing system (10) as claimed in Claim 15, wherein said distributed cache memory (14, 16) further includes a memory controller cache (48) within said memory controller (20).

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FIG. 1

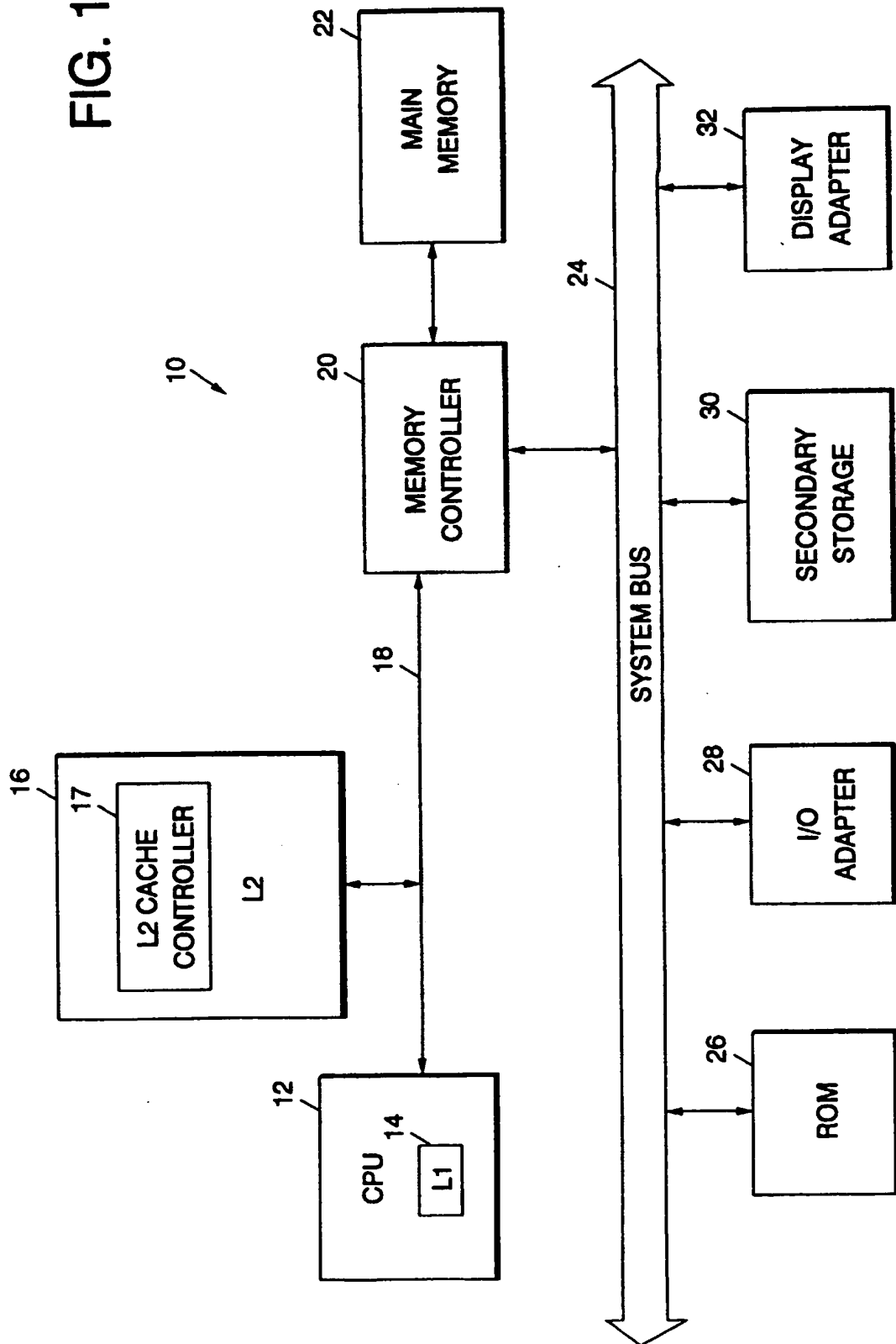


FIG. 2

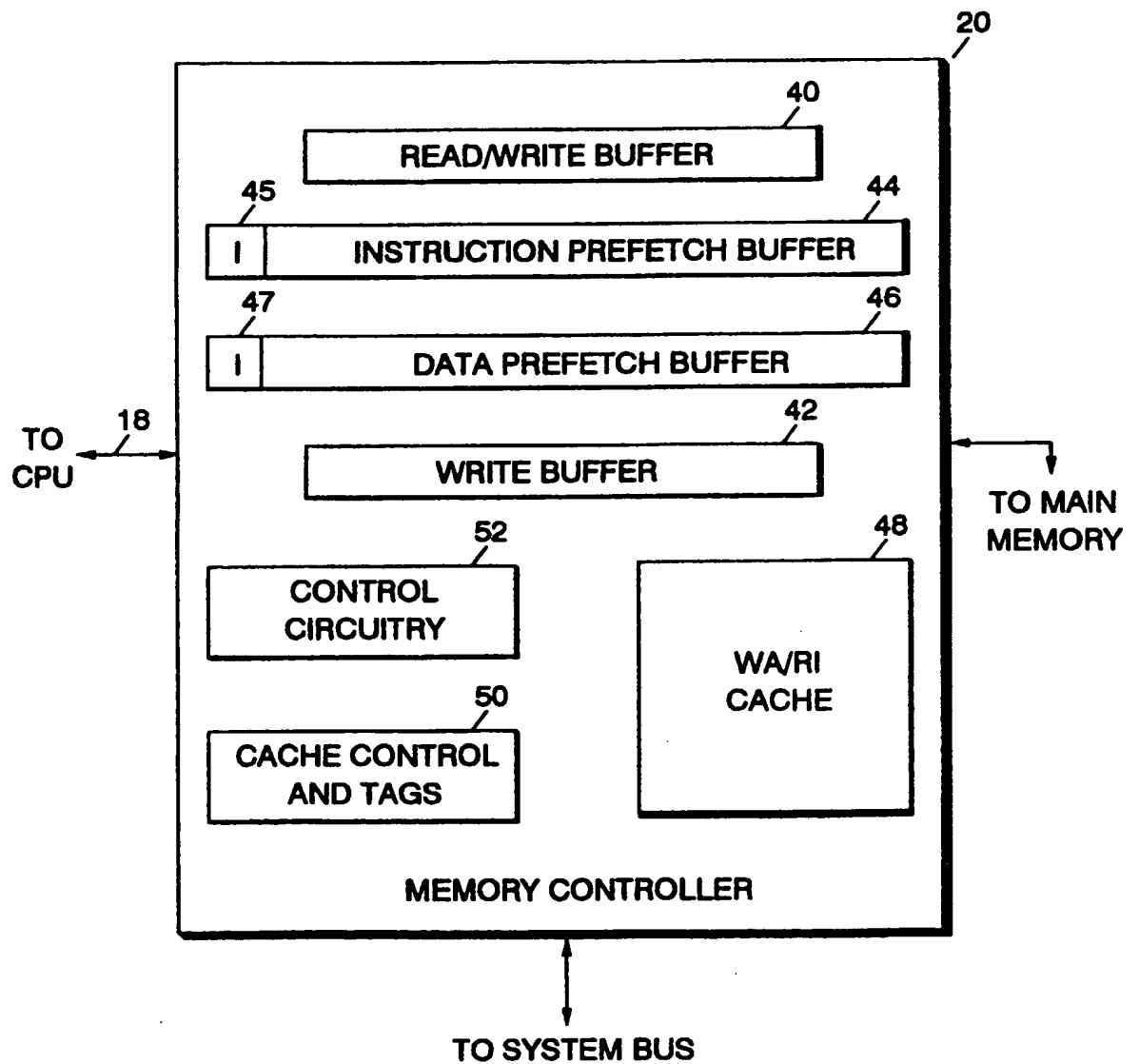
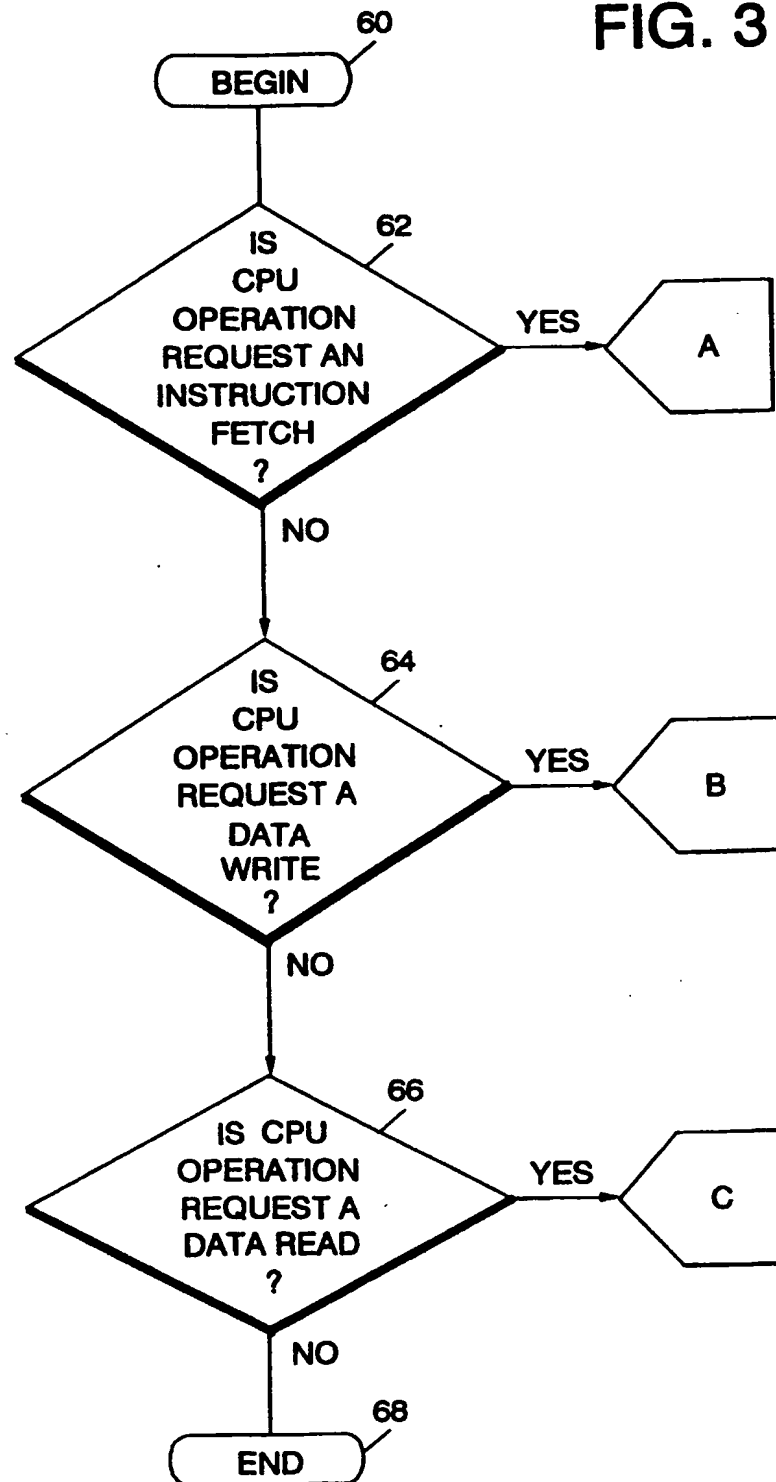


FIG. 3



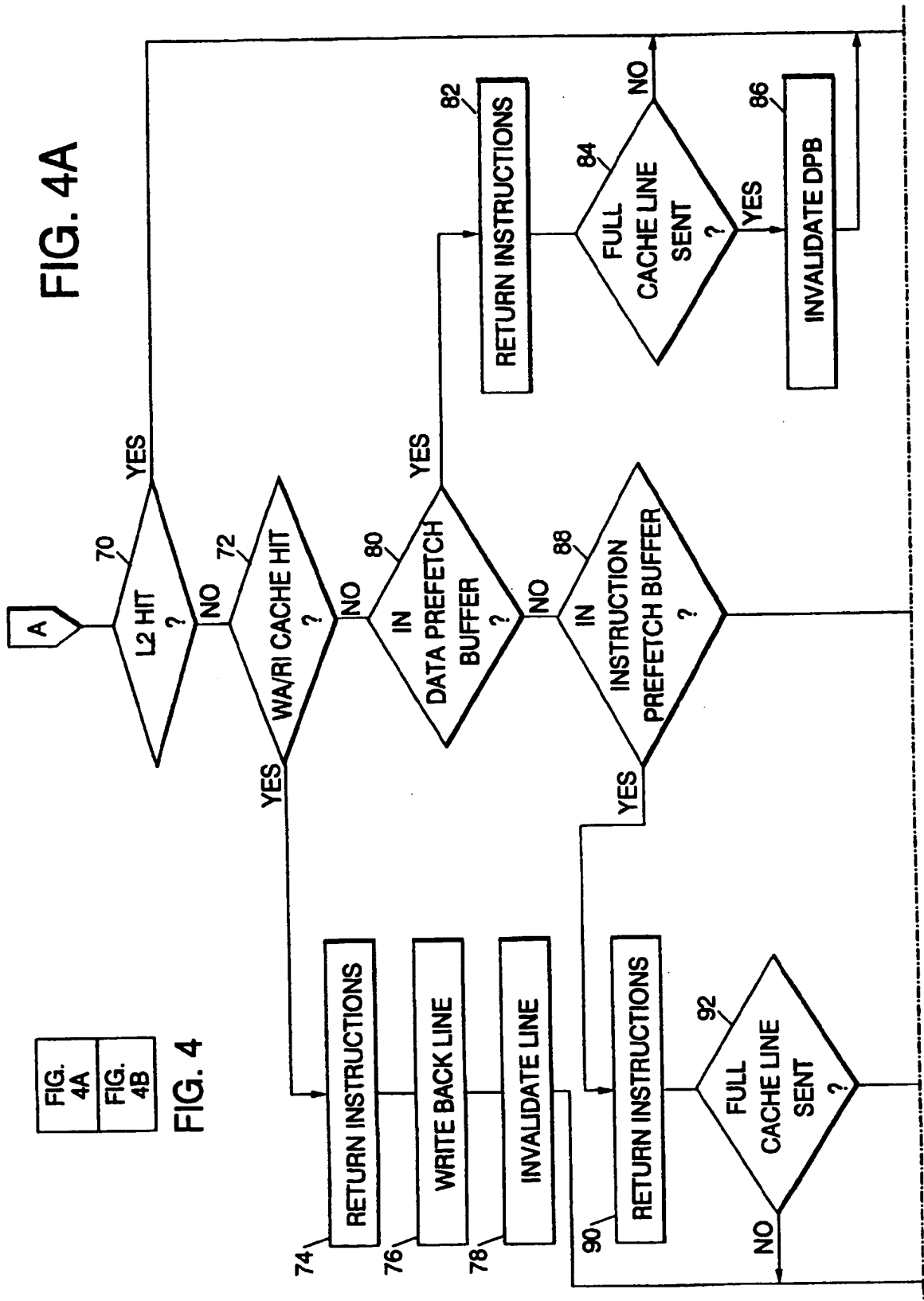


FIG. 4A  
FIG. 4B

FIG. 4

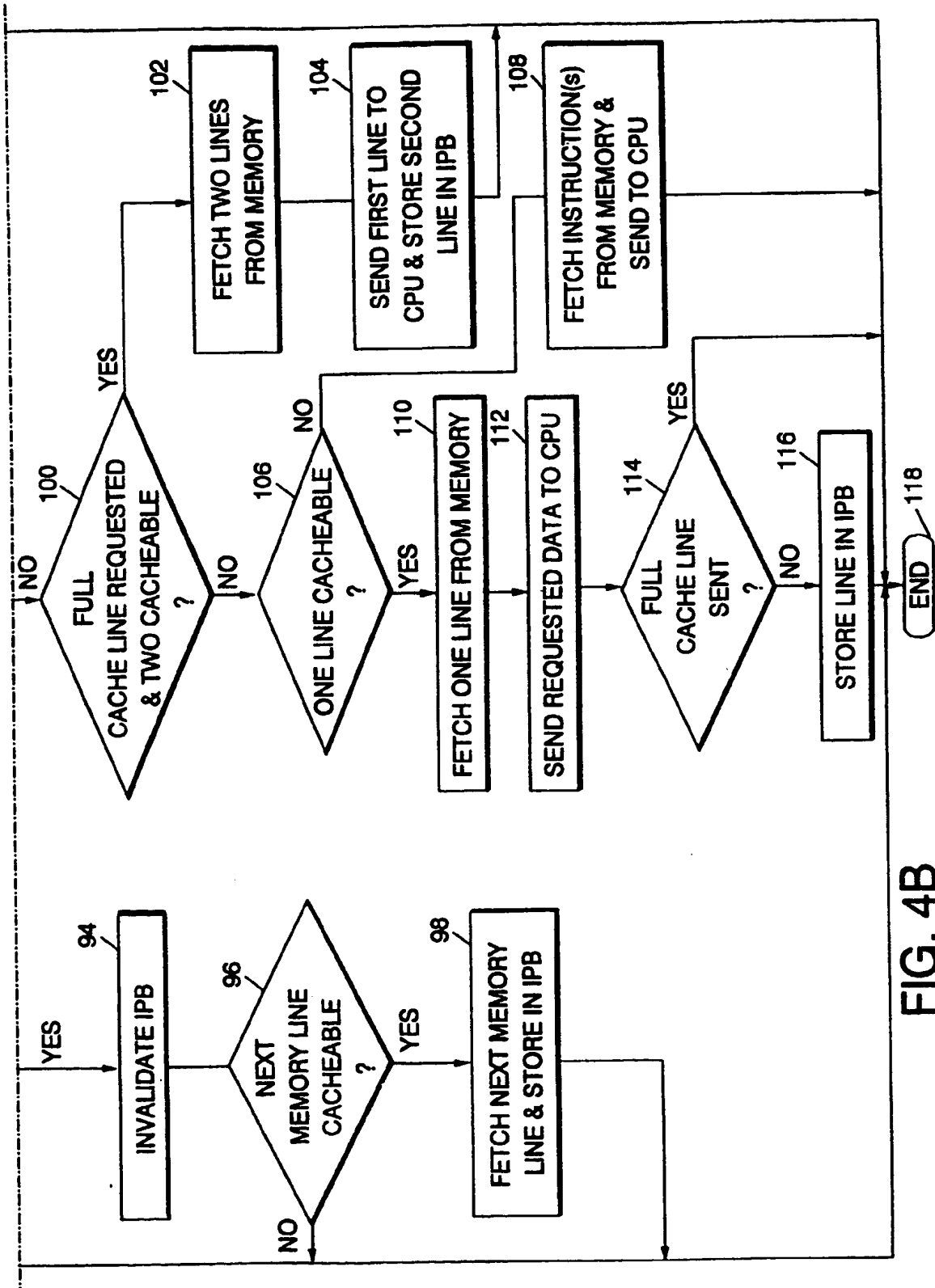
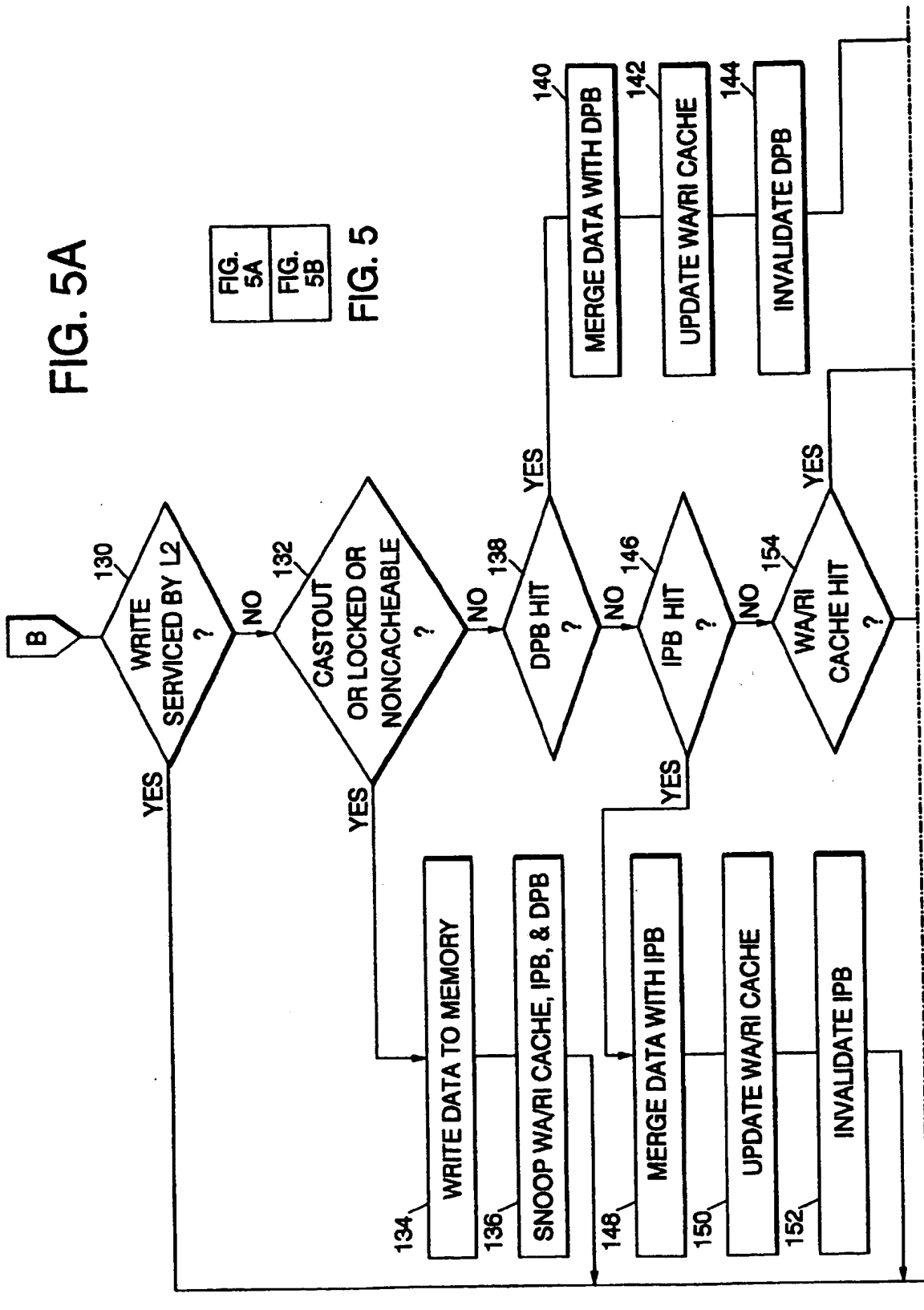
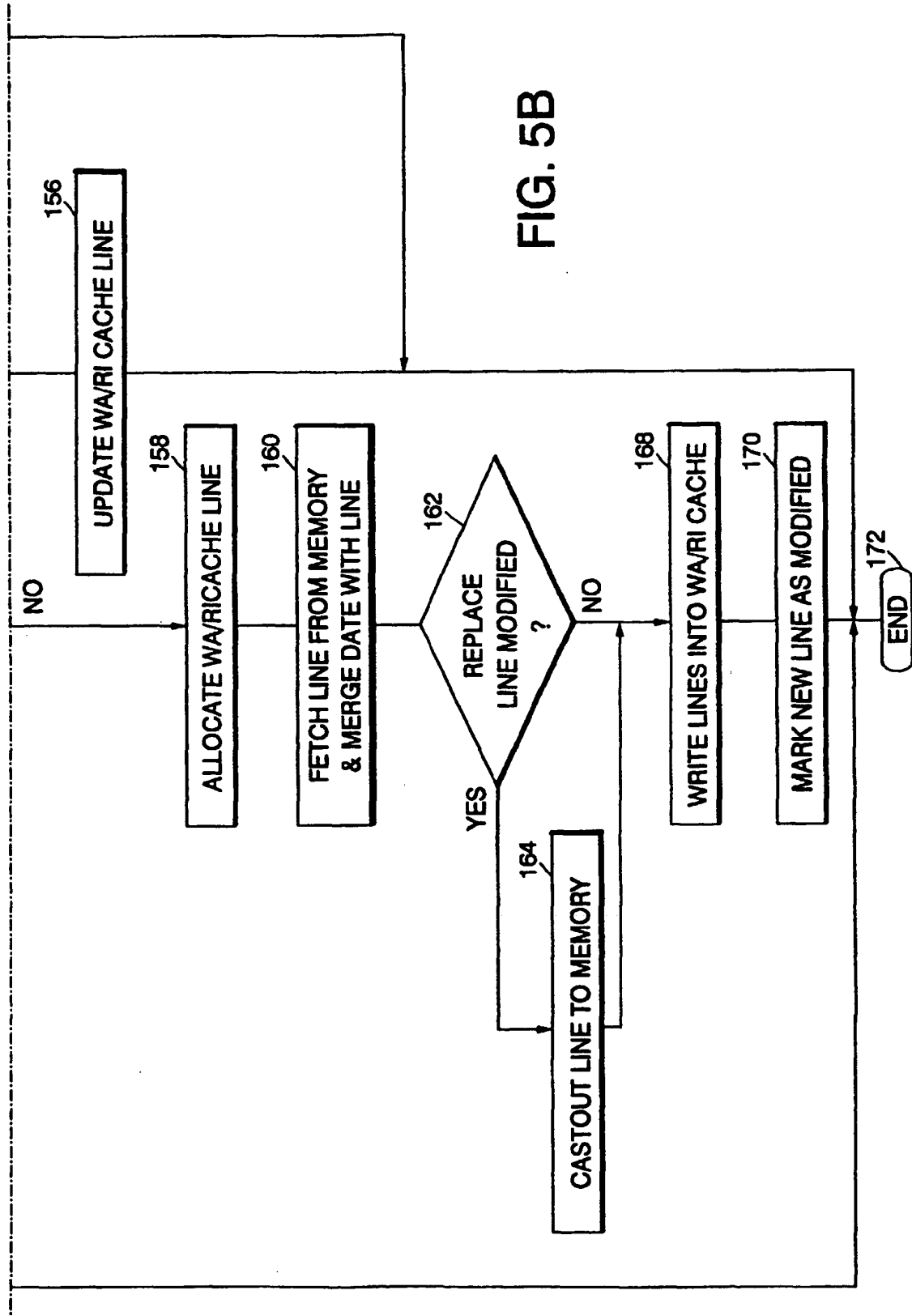


FIG. 4B







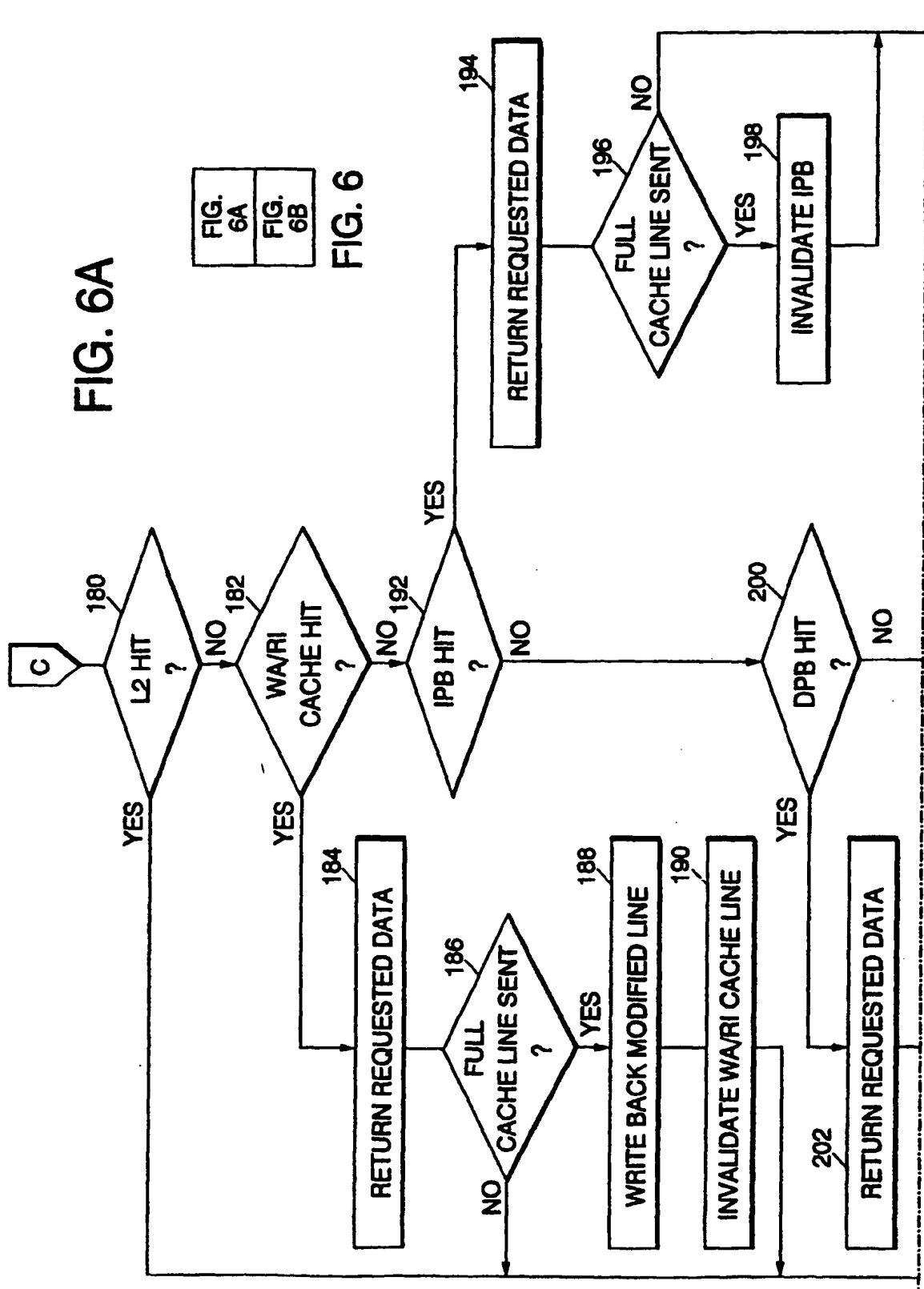


FIG. 6A  
FIG. 6B  
FIG. 6

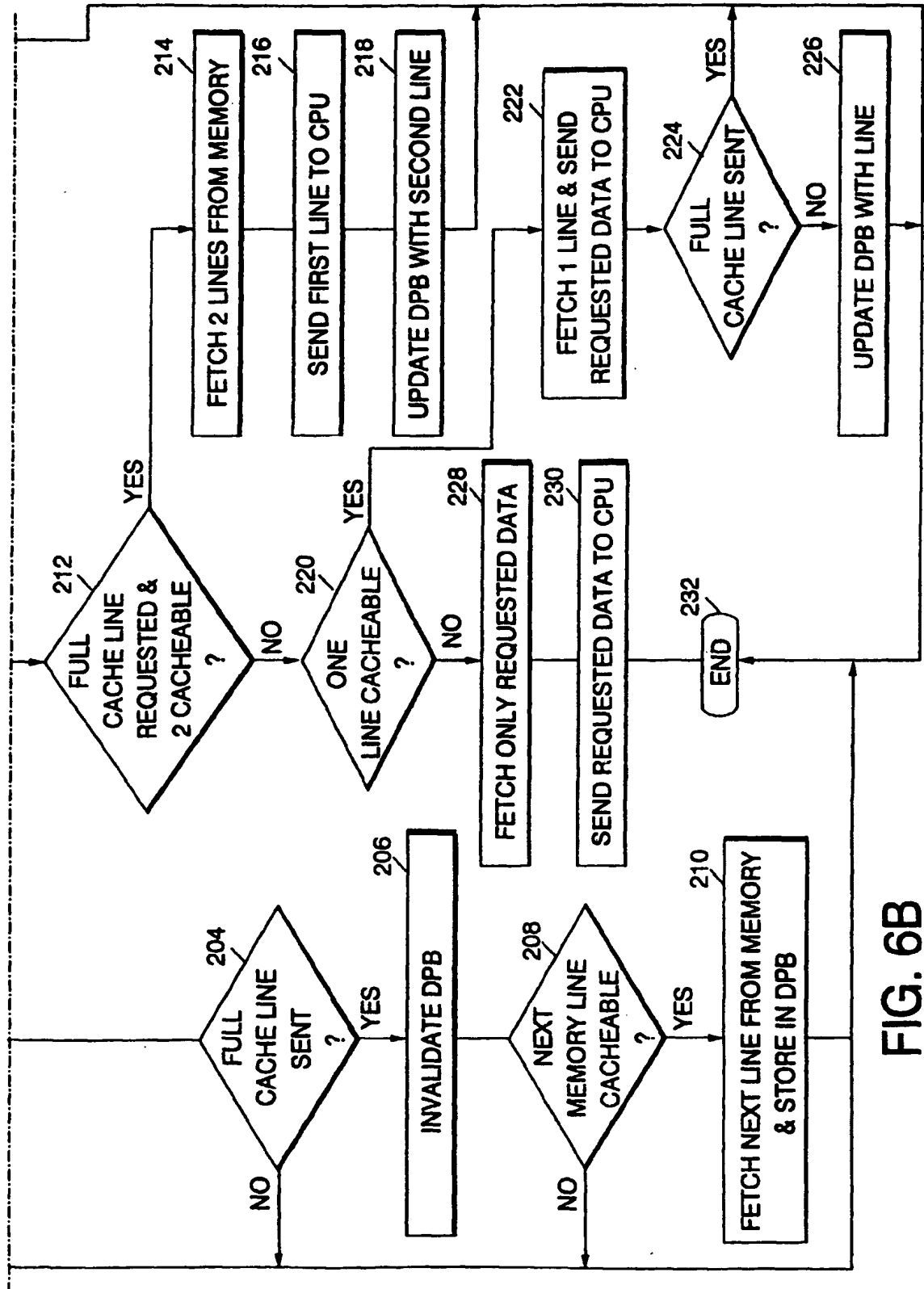


FIG. 6B



European Patent  
Office

# EUROPEAN SEARCH REPORT

Application Number  
EP 97 30 1738

DOCUMENTS CONSIDERED TO BE RELEVANT			
Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int.Cl.6)
X A	US 5 367 657 A (KHARE MANOJ ET AL) 22 November 1994 * column 3, line 1 - column 5, line 36; figures 1,2 *	5,6,14, 15 7-13,16	G06F12/08
X A	EP 0 604 139 A (NCR INT INC) 29 June 1994 * column 4, line 29 - line 54; figures 1-3 *	5,6,14 7-13,15, 16	
X Y	EP 0 470 735 A (NCR CO) 12 February 1992 * column 2, line 34 - column 5, line 3; figure 1 *	5-7,9, 10,14-16 1-4	
Y	EP 0 468 831 A (DIGITAL EQUIPMENT CORP) 29 January 1992 * page 39, line 32 - page 42, line 34; figures 19,22,23 *	1-4	
A	GB 2 276 964 A (HYPERTECH PTY LTD) 12 October 1994 * page 2, line 6 - page 3, line 24; figures 2,6 *	1-4	TECHNICAL FIELDS SEARCHED (Int.Cl.6)
A	WO 93 18459 A (RAMBUS INC) 16 September 1993 * page 3, line 1 - page 4, line 10; figures 1,3 *	5-10	G06F
The present search report has been drawn up for all claims			
Place of search THE HAGUE		Date of completion of the search 7 July 1997	Examiner Nielsen, O
<p><b>CATEGORY OF CITED DOCUMENTS</b></p> <p>X : particularly relevant if taken alone Y : particularly relevant if combined with another document of the same category A : technological background O : non-written disclosure P : intermediate document</p> <p>T : theory or principle underlying the invention E : earlier patent document, but published on, or after the filing date D : document cited in the application L : document cited for other reasons &amp; : member of the same patent family, corresponding document</p>			